Transactions

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Transaction

 sequence of one or more SQL statements to be treated as a one atomic unit
 carried out in isolation

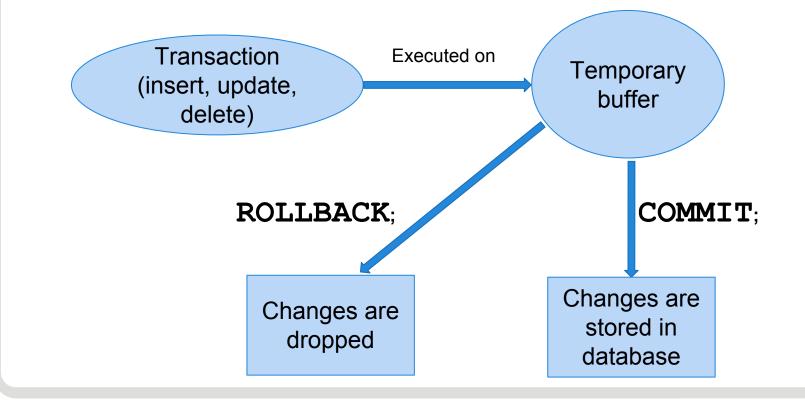
 no operation may be performed during

this transaction to affect it

 at time of system failure either all changes are applied or no changes are applied at all
 DB remains in a consistent state

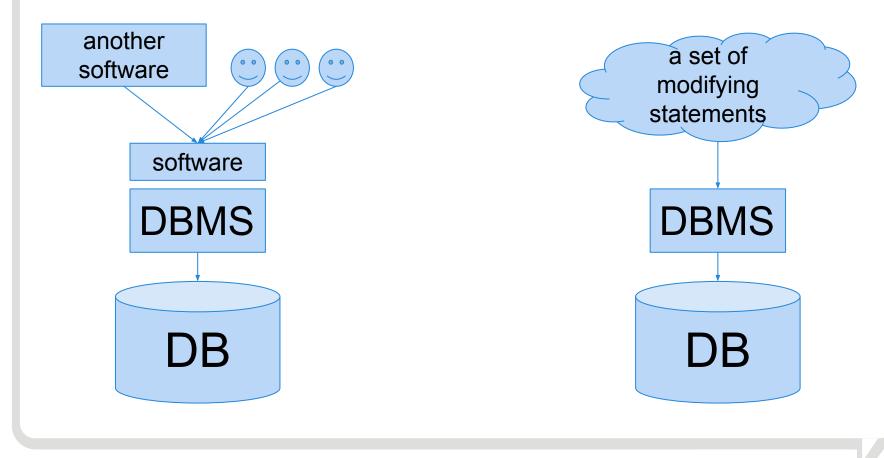
Transactional execution

 Sequence of modifying statements executed as a whole - transaction
 O Either all changes are stored
 Or none is stored



Motivation for transaction

- Simultaneous (parallel) access to the database
- Resistance to system failures



Goals of transactions

- parallel access to the database
 - executing a sequence of commands as if performed in isolation
 - allows simultaneous execution whenever possible
- resistance to system failures
 - guarantee that either everything is executed or nothing is executed, regardless of possible system failure

Parallel access

- database system should support parallel access for multiple users
- while maintaining the integrity and consistency of data
- during paralell access different types of inconsistency can occur:
 - $\circ\,\text{at}$ the attribute level
 - \circ at the record (row) level
 - $\circ\,\text{at}$ the table level
 - $\circ\,\text{at}$ the level of multiple statements

Example

UPDATE account
SET balance = balance - 100
WHERE id = 21;
UPDATE account
SET balance = balance + 100
WHERE id = 22;

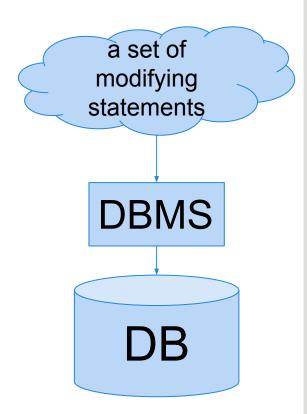
UPDATE account

SET balance = balance * 1.02

WHERE type = 'common';

Resistance to system failures

- the system should ensure data consistency even in case of system failure
- failure during a series of commands (database initialization)
- during data migration in DB
 while the updates are in a memore (not stored to hard drive, yet)



Example

UPDATE account SET balance = balance - 100 WHERE id = 21;

UPDATE account SET balance = balance + 100 WHERE id = 22;

Transactions in SQL

- starts automatically with the first statement
- using the commit keyword command will finish the transaction and start a new one
- the current transaction ends at the end of the session
- autocommit mode: each statement is in it's own transaction
- changes done in current transaction can be canceled by the rollback keyword command

Example

```
START TRANSACTION;
INSERT INTO student
VALUES (10, 'Janko', 'Hraško', 4);
COMMIT;
```

```
START TRANSACTION;
UPDATE student
SET name = 'Jozko'
WHERE id = 10;
ROLLBACK;
```

ACID

atomicity

o all or nothing (rollback allows to cancel
 all changes made in transaction)

consistency

 for each transaction, it can be assumed that all DB restrictions are met before it starts and must ensure that are met after it is completed

ACID

• isolation

- serializability the result of operations corresponds to some sequential execution of all transactions
- 4 levels of isolation (Serializable, Repeatable Read, Read Committed, Read Uncommitted)
- durability
 - after commit, all changes must remain in the database

